

Higher-Order Functions (Pre Lecture)

Dr. Neil T. Dantam

CSCI-400, Colorado School of Mines

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Introduction

Higher-Order Functions

- ▶ **Higher-order function:** takes or returns another function
- ▶ The common ones:
 - map transforms elements
 - filter eliminates elements
 - fold combines elements
- ▶ Useful abstractions in functional and non-functional settings

Outcomes

- ▶ Understand the operation of map, filter, and fold
- ▶ Apply map, filter, and fold to operate on collections

Outline

Common Higher-Order Functions

- Map

- Filter

- Fold (aka Reduce)

Usage

Higher-order functions

Definition: Higher-order function

A function that takes another function as an argument or returns another function as its result.

Example (Passing)

Function $f(g,a)$

1 `return g(42, a)`

Example (Returning)

Function $f(a)$

```
1 function g(b) is
2   |   return a + b
3 return g
```

Counterexample

Function $f(a,b)$

1 `return a + b`

Map function

Definition (map)

Apply a function to every member of an input sequence, and collect the results into the output sequence.

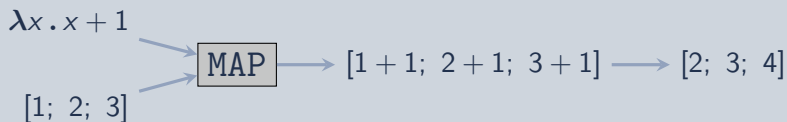
$$\text{map} : \underbrace{(\mathbb{X} \mapsto \mathbb{Y})}_{\text{function}} \times \underbrace{\mathbb{X}^n}_{\text{input sequence}} \mapsto \underbrace{\mathbb{Y}^n}_{\text{output sequence}}$$

Illustration



Example: Map

Increment every element by one



Not ($\neg$) every element



Algorithm: Map function on lists

Functional Map

Function $\text{map}(f, \ell)$

```

1 match  $\ell$  with
2   | case NIL  $\rightarrow$  NIL
3   | case  $\text{cons}(h, t) \rightarrow$ 
4     |  $\text{cons}(f(h), \text{map}(f, t))$ 

```

Imperative Map

Procedure $\text{map}(f, \ell)$

```

1  $Y \leftarrow \text{make-sequence}(|\ell|)$ 
2  $i \leftarrow 0$ 
3 while  $i < n$  do
4   |  $Y[i] \leftarrow f(\ell[i])$ 
5   |  $i \leftarrow i + 1$ 
6 return  $Y$ 

```

Exercise 1: Map

Function `double-list( $\ell$ )`

Filter

Definition (filter)

Apply a predicate (Boolean-valued function) to every member of an input sequence, and collect the only the input elements where the predicate is true.

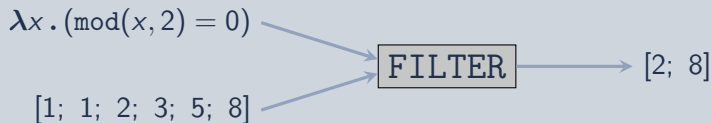
$$\text{filter} : \underbrace{(\mathbb{X} \mapsto \mathbb{B})}_{\text{function}} \times \underbrace{\mathbb{X}^m}_{\text{input sequence}} \mapsto \underbrace{\mathbb{X}^n}_{\text{output sequence}}$$

Illustration



Example: Filter

Evens



Odds



Algorithm: Filter

Filter Algorithm

Function $\text{filter}(p, s)$

```

1 match  $s$  with
2   | case NIL  $\rightarrow$  NIL
3   | case  $\text{cons}(f, r) \rightarrow$ 
4     | if  $p(f)$  then
5     |   |  $\text{cons}(f, \text{filter}(p, r))$ 
6     | else
7     |   |  $\text{filter}(p, r)$ 

```

Exercise 2: Filter

Function `remove-empty( $\ell$ )`

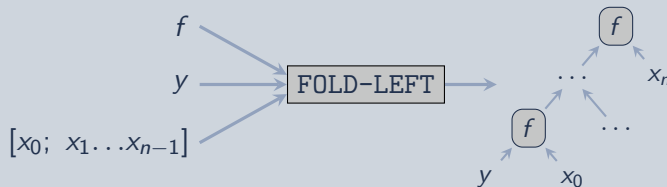
Fold-left

Definition (fold-left)

Apply a binary function to each member of a sequence and the prior result, starting from the left.

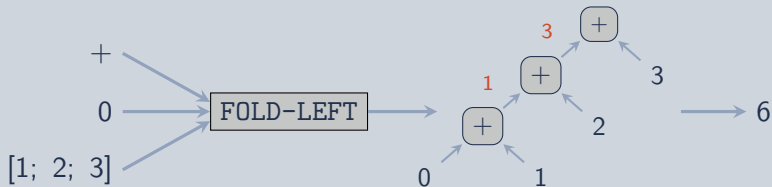
$$\text{fold-left} : \underbrace{(\mathbb{Y} \times \mathbb{X} \mapsto \mathbb{Y})}_{\text{function}} \times \underbrace{\mathbb{Y}}_{\text{init.}} \times \underbrace{\mathbb{X}^n}_{\text{sequence}} \mapsto \underbrace{\mathbb{Y}}_{\text{result}}$$

Illustration

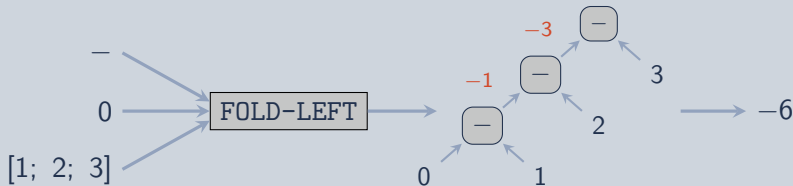


Example: Fold-left

Example (Addition)



Example (Subtraction)



Algorithm: Fold-left

Functional

Function fold-left(f , y , X)

```

1 match  $X$  with
2   | case NIL  $\rightarrow y$ 
3   | case cons( $a$ ,  $X'$ )  $\rightarrow$ 
4     | fold-left( $f$ ,  $f(y, a)$ ,  $X'$ )

```

Imperative

Procedure fold-left(f , y , X)

```

1  $i \leftarrow 0$ 
2 while  $i < |X|$  do
3   |  $y \leftarrow f(y, X)$ 
4   |  $i \leftarrow i + 1$ 
5 return  $y$ 

```

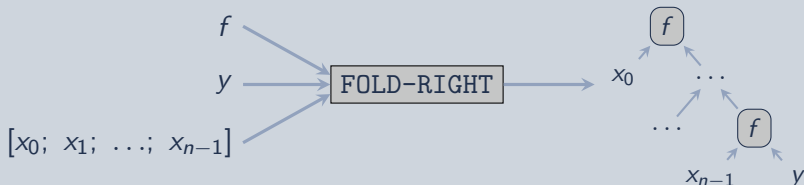
Fold-right

Definition (fold-right)

Apply a binary function to each member of a sequence and the prior result, starting from the right.

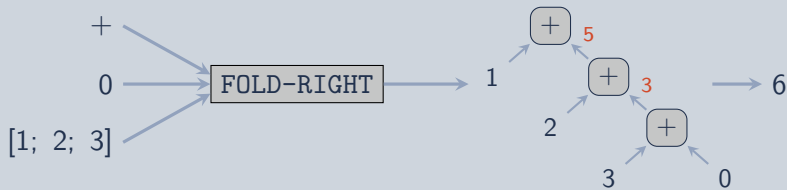
$$\text{fold-right} : \underbrace{(\mathbb{X} \times \mathbb{Y} \mapsto \mathbb{Y})}_{\text{function}} \times \underbrace{\mathbb{Y}}_{\text{init.}} \times \underbrace{\mathbb{X}^n}_{\text{sequence}} \mapsto \underbrace{\mathbb{Y}}_{\text{result}}$$

Illustration

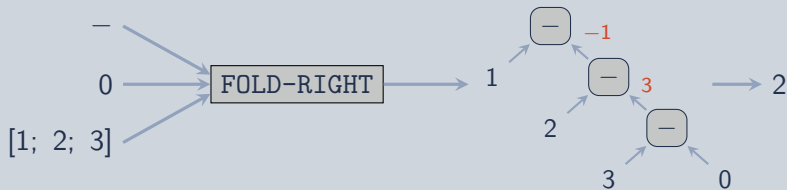


Example: Fold-right

Example (Addition)



Example (Subtraction)



Algorithm: Fold-right

Recursive

Function fold-right(f , y , X)

```

1 match  $X$  with
2   | case NIL  $\rightarrow y$ 
3   | case cons( $a$ ,  $X'$ )  $\rightarrow$ 
4     |  $f(a, \text{fold-right}(f, y, X'))$ 

```

Procedural

Procedure fold-right(f , y , X)

```

1  $i \leftarrow |X| - 1$ 
2 while  $i \geq 0$  do
3   |  $y \leftarrow f(X_i, y)$ 
4   |  $i \leftarrow i - 1$ 
5 return  $y$ 

```

Outline

Common Higher-Order Functions

Map

Filter

Fold (aka Reduce)

Usage

Generalized Operations

- ▶ `map`, `filter`, `fold` are often defined on lists
- ▶ Possible to define on other collections or data types:
 - ▶ Arrays
 - ▶ Trees
 - ▶ Sets
 - ▶ key-value pairs (aka “maps”)

Fold is Fascinating

Loop

Procedure $f(\ell)$

```

1  $a \leftarrow \text{init}$ 
2 while  $\neg \text{empty}(\ell)$  do
3    $f \leftarrow \text{first}(\ell)$ 
4    $\ell \leftarrow \text{rest}(\ell)$ 
5    $a \leftarrow \text{frob}(a, f)$ 
6 return  $a$ 

```

Recursion

Function $f(\ell)$

```

1 function  $h(\ell, a)$  is
2   match  $\ell$  with
3     case  $\text{NIL}$   $\rightarrow a$ 
4     case  $\text{cons}(f, r)$   $\rightarrow$ 
5        $h(r, \text{frob}(a, f))$ 
6  $h(\ell, \text{init})$ 

```

Fold

Function $f(\ell)$

```

1  $\text{fold}(\text{frob}, \text{init}, \ell)$ 

```

Compact, Abstract, General

Example: Length via Fold

Tail-Recursive

Function $\text{length}(\ell)$

```

1 function  $h(\ell, y)$  is
2   match  $\ell$  with
3   | case  $\text{NIL} \rightarrow y$ 
4   | case  $\text{cons}(x, d) \rightarrow$ 
5   |    $h(d, 1 + y)$ 
6  $h(\ell, 0)$ 

```

Fold

Function $\text{length}(\ell)$

```

1 fold-left( $\lambda y x. 1 + y$ , 0,  $\ell$ )

```

Example: Length via Fold

`length([1; 2; 3])`

fold-left

Function `fold-left( $f, y, \ell$ )`

```

1 match  $\ell$  with
2   | case NIL  $\rightarrow y$ 
3   | case cons( $x, r$ )  $\rightarrow$ 
4     | fold-left( $f, f(y, x), r$ )

```

length

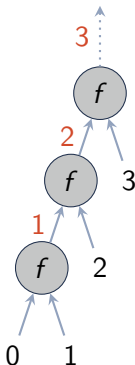
Function `length( $\ell$ )`

```

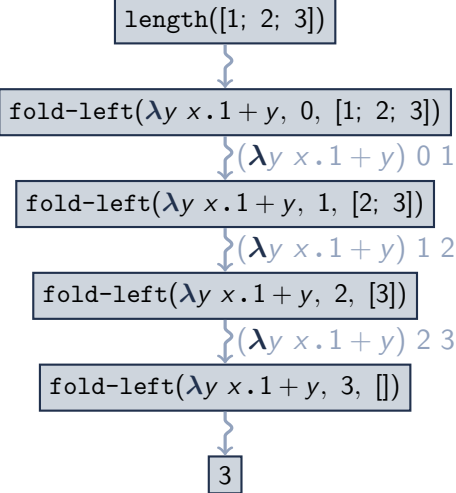
1 fold-left( $\lambda y x. 1 + y, 0, \ell$ )

```

Calls



Trace



Using Fold

Questions

- ▶ What is the “result” (y)?
- ▶ How do I use the current element (x) to update the the result?
- ▶ Do I need to start from the beginning or end of the list?

Fold-Left

```

      result  current element
function h( $\overbrace{y}$ ,  $\overbrace{x}$ ) is
  | /* Use  $x$  to update  $y$  */
  | ...
fold-left(h, initial-value,  $\ell$ )
  
```

Fold-Right

```

      current element  result
function h( $\overbrace{x}$ ,  $\overbrace{y}$ ) is
  | /* Use  $x$  to update  $y$  */
  | ...
fold-right(h, initial-value,  $\ell$ )
  
```


Exercise 3: Count via Fold

Tail-Recursive

Function $\text{count}(\ell, p)$

```

1 function  $h(\ell, a)$  is
2   match  $\ell$  with
3     case NIL  $\rightarrow a$ 
4     case cons  $(f, r) \rightarrow$ 
5       if  $p(f)$  then
6         |  $h(r, 1 + a)$ 
7       else
8         |  $h(r, a)$ 
9  $h(\ell, 0)$ 

```

Fold

Function $\text{count}(\ell, p)$

Example: Copy via Fold

Recursive

Function `copy( $\ell$ )`

```
1 match  $\ell$  with
2   | case NIL  $\rightarrow$  NIL
3   | case cons( $a$ ,  $d$ )  $\rightarrow$ 
4     |   cons( $a$ , copy( $d$ ))
```

Fold

Function `copy( $\ell$ )`

```
1 fold-right(cons, NIL,  $\ell$ )
```

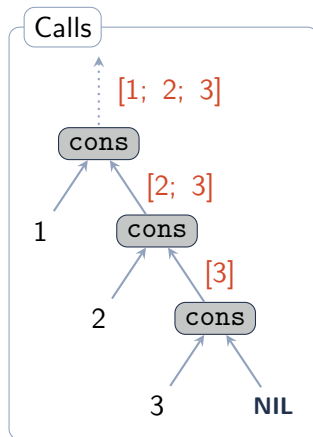
Example: Copy via Fold

`copy([1; 2; 3])`

`copy`

Function `copy( $\ell$ )`

1 `fold-right(cons, NIL,  $\ell$ )`



Exercise 4: Append via Fold

Recursive

Function $\text{append}(\ell_a, \ell_b)$

```
1 match  $\ell_a$  with
2   | case NIL  $\rightarrow \ell_b$ 
3   | case  $\text{cons}(a, d) \rightarrow$ 
4     |  $\text{cons}(a, \text{append}(d, \ell_b))$ 
```

Recursive

Function $\text{append}(\ell_a, \ell_b)$

Exercise 5: Reverse via Fold

Recursive

Function $\text{reverse}(\ell)$

```

1 function  $h(\ell, r)$  is
2   match  $\ell$  with
3   | case  $\text{NIL} \rightarrow r$ 
4   | case  $\text{cons}(a, d) \rightarrow$ 
5   |    $h(d, \text{cons}(a, r))$ 
6  $h(\ell, \text{NIL})$ 

```

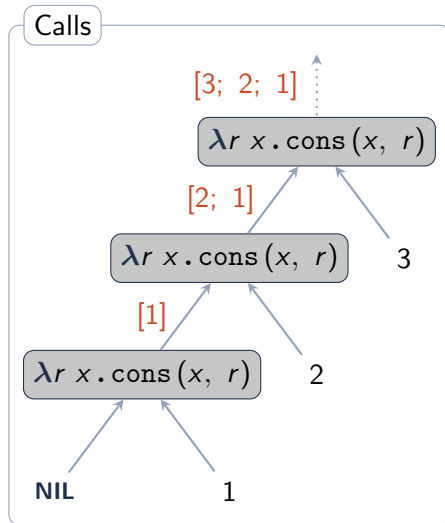
Fold

Function $\text{reverse}(\ell)$

Exercise 5: Reverse via Fold

`reverse([1; 2; 3])`

`reverse`



Example: Map via Fold

Recursive Map

Function $\text{map}(f, \ell)$

```

1 match  $\ell$  with
2   | case NIL  $\rightarrow$  NIL
3   | case  $\text{cons}(x, r) \rightarrow$ 
4     |  $\text{cons}(f(x), \text{map}(f, r))$ 

```

Copy via Fold

Function $\text{copy}(\ell)$

```

1 fold-right( $\lambda x y. \text{cons}(x, y)$ , NIL,  $\ell$ )

```

Map via Fold

Function $\text{map}(f, \ell)$

```

1 fold-right( $\lambda x y. \text{cons}(f(x), y)$ , NIL,  $\ell$ )

```

Example: Map via Fold

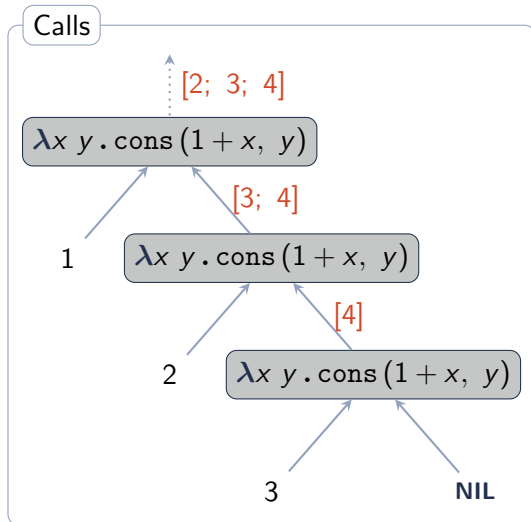
$\text{map}(\lambda x. 1 + x, [1; 2; 3])$

map

Function $\text{map}(f, \ell)$

```

1 function  $h(x, y)$  is
2   |  $\text{cons}(f(x), y)$ 
3 fold-right( $h$ , NIL,  $\ell$ )
  
```



Exercise 6: Filter via Fold

Recursive

Function $\text{filter}(p, \ell)$

```
1 match  $\ell$  with
2   | case NIL  $\rightarrow$  NIL
3   | case cons( $x, r$ )  $\rightarrow$ 
4     | if  $p(x)$  then
5       |   cons( $x, \text{filter}(p, r)$ )
6       | else
7       |   filter( $p, r$ )
```

Fold

Function $\text{filter}(p, \ell)$

Every

Definition (every)

Return true when predicate evaluates to true on every member of the input sequence. Otherwise, return false.

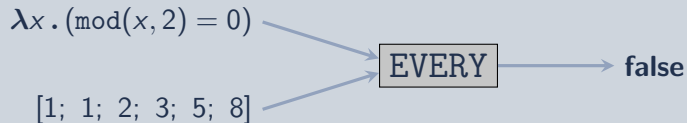
$$\text{every} : \underbrace{(\mathbb{X} \mapsto \mathbb{B})}_{\text{function}} \times \underbrace{\mathbb{X}^n}_{\text{input sequence}} \mapsto \mathbb{B}$$

Illustration

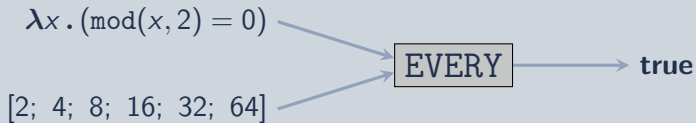


Example: Every

Evens



Evens



Exercise 7: Every via Fold

Recursive

Function $\text{every}(p, \ell)$

```
1 match  $\ell$  with
2   | case NIL  $\rightarrow$  true
3   | case cons( $x, r$ )  $\rightarrow$ 
4     | if  $p(x)$  then
5       |   every( $p, r$ )
6     | else
7       |   false
```

Fold

Function $\text{every}(p, \ell)$

```
1 function  $h(y, x)$  is
2   | if  $p(x)$  then  $y$ 
3   | else false
4 fold-left( $h, \text{true}, \ell$ )
```

What's a potential issue?

Efficiency Notes

- ▶ Fold-left is tail recursive:
any decent implementation on lists will use constant $O(1)$ space
- ▶ Fold-right is *not* tail recursive:
an implementation on lists will (likely) use linear $O(|\ell|)$ stack space
- ▶ Folding over other structures (arrays, trees) may need different amounts of space.
- ▶ Fold must visit every element, so time is linear $O(|\ell|)$

Application: MapReduce

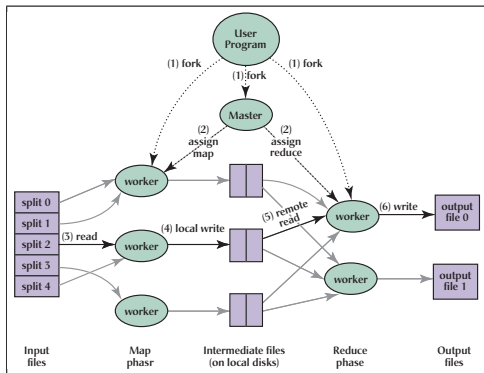
Function $\text{MapReduce}(f, g, X)$

```

1  $Y \leftarrow \text{parallel-map}(f, X)$ 
2 return  $\text{reduce}(g, Y)$ 

```

- ▶ Idea:
 - ▶ (parallel) map
 - ▶ (serial) reduce/fold
- ▶ Provides scalability, fault-tolerance
- ▶ Implementations:
 - ▶ Google MapReduce
 - ▶ Apache Hadoop



Jeffrey Dean and Sanjay Ghemawat.

MapReduce: Simplified Data Processing on Large Clusters. Communications of the ACM. 2008.

Summary

- ▶ Common Higher-order Functions:
 - `map` transforms elements
 - `filter` eliminates elements
 - `fold` combines elements
- ▶ Compact and Abstract:
 - ▶ Factors out control to operate over a collection
 - ▶ Per-element operation is independent of collection type

References

- ▶ Clarkson. Functional Programming in OCaml.
 - ▶ Ch 4 Higher-order programming